



Well-quasi-ordering, clique-width and factorial properties of graphs

Vadim Lozin

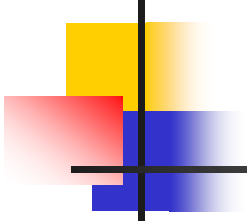
DIMAP – Center for Discrete Mathematics and its Applications

Mathematics Institute

University of Warwick

Is there anybody going to listen to my story?

The Beatles, "Girl"



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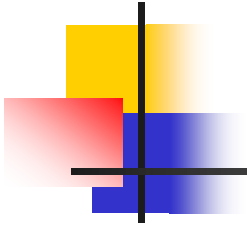
Definition. *A class of graphs is hereditary if it is closed under taking induced subgraphs*

Observation. *A class of graphs is hereditary if and only if it can be characterized in terms of forbidden induced subgraphs.*

Bipartite graphs = $\text{Free}(C_3, C_5, C_7, \dots)$: graphs partitionable into two independent sets

Split graphs = $(C_4, C_5, 2K_2)$: graphs partitionable into a clique and an independent set

The speed of hereditary properties of graphs

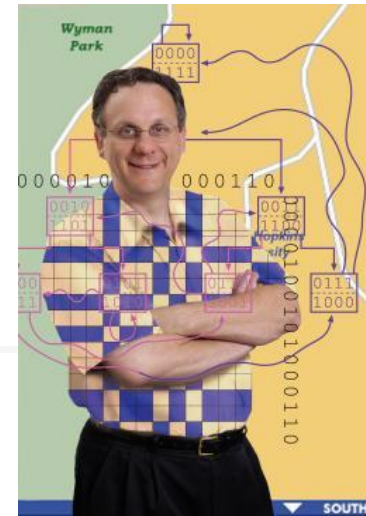




The speed of hereditary properties of graphs

(the number of n -vertex labeled graphs)

The speed of hereditary properties of graphs.



Edward R. Scheinerman

E.R. Scheinerman, J. Zito,
On the size of hereditary
classes of graphs. *J. Combin.*
***Theory Ser. B* 61 (1994), no.**
1, 16–39.

Factorial

Exponential

Polynomial

Constant

The speed of hereditary properties of graphs.



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Issled. Oper. Ser. 1 4 (1997) 3–12.



J. Balogh, B. Bollobás, D. Weinreich,

The speed of hereditary properties of graphs.

J. Combin. Theory Ser. B 79 (2000) 131–156.



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Factorial classes of graphs

- interval graphs
- permutation graphs
- line graphs
- threshold graphs
- cographs
- forests
- planar graphs
- all proper minor-closed graph classes
- all classes of graphs of bounded vertex degree



Conjecture

A class X of graphs is factorial if and only if the fastest of the following three classes is factorial:

- the class of bipartite graphs in X
- the class of co-bipartite graphs in X
- the class of split graphs in X

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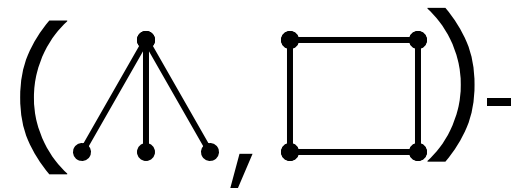


Graph classes defined by forbidden induced subgraphs with at most 4 vertices

V. Lozin, C. Mayhill, V. Zamaraev **A note on the speed of hereditary properties of graphs.** *Electronic J. Combinatorics*, 18 (1) 2011, #P157.

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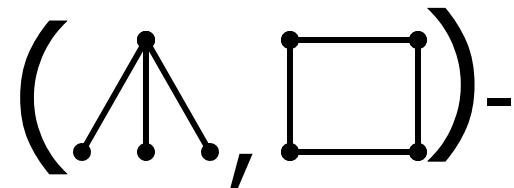


free graphs

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[V. Lozin, C. Mayhill, V. Zamaraev](#) **A note on the speed of hereditary properties of graphs.** *Electronic J. Combinatorics*, 18 (1) 2011, #P157.

[Chudnovsky, Maria](#); [Seymour, Paul](#) **The structure of claw-free graphs.** *Surveys in combinatorics 2005*, 153–171, [London Math. Soc. Lecture Note Ser., 327](#), Cambridge Univ. Press, Cambridge, 2005.



free graphs



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"Smallest" superfactorial class of bipartite graphs



The diagram is a funnel shape, wider at the top and narrower at the bottom. It is divided into five horizontal sections by four lines. The top section is the largest and contains the text "C₄-free bipartite graphs". The subsequent sections, from top to bottom, are labeled "Factorial", "Exponential", "Polynomial", and "Constant". The funnel narrows as it descends, indicating that each subsequent class is a smaller subset of the one above it.

C_4 -free bipartite graphs

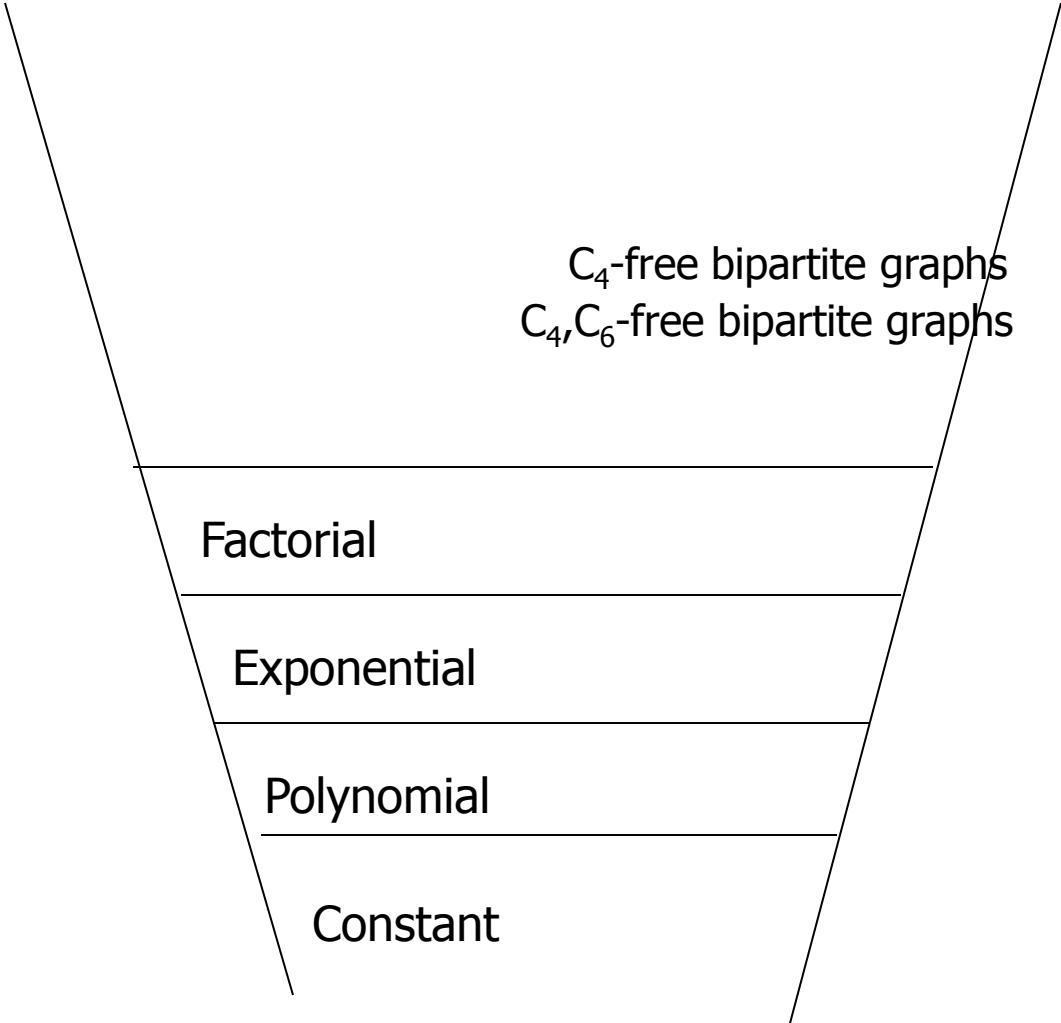
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C_4 -free bipartite graphs
 C_4, C_6 -free bipartite graphs

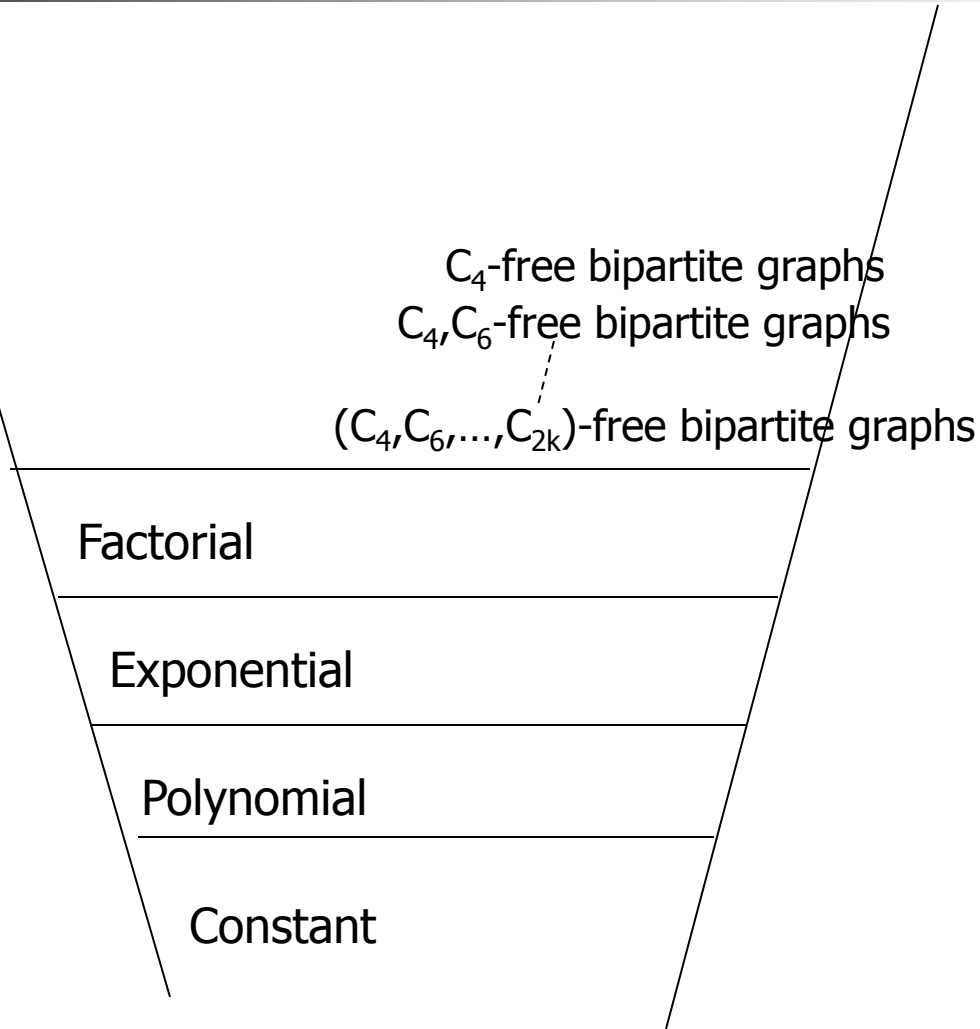
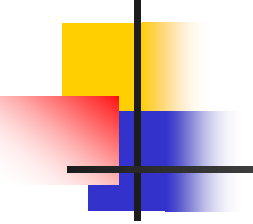
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"Smallest" superfactorial class of bipartite graphs

C_4 -free bipartite graphs
 C_4, C_6 -free bipartite graphs
 $(C_4, C_6, \dots, C_{2k})$ -free bipartite graphs

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Lazebnik, F.; Ustimenko, V. A.;
Woldar, A. J. **A new series of
dense graphs of high girth.**
Bull. Amer. Math. Soc. (N.S.) 32
(1995), no. 1, 73–79.

"Smallest" superfactorial class of bipartite graphs

Chordal bipartite graphs C_4 -free bipartite graphs
Free($C_3, C_5, C_6, C_7, \dots$) C_4, C_6 -free bipartite graphs
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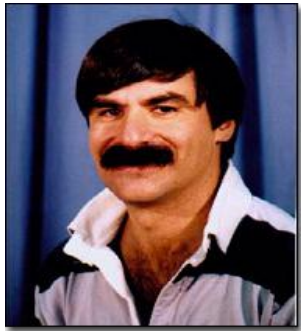
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Spinrad, Jeremy P.
Nonredundant 1's
in Γ -free matrices.
SIAM J. Discrete
Math. 8 (1995), no. 2,
 251–257



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- all classes of graphs of bounded clique-width



Any hereditary class of graphs of bounded clique-width is at most factorial

P. Allen, V. Lozin, M. Rao **Clique-width and the speed of hereditary properties.** *Electronic J. Combinatorics*, 16 (2009), #R35.

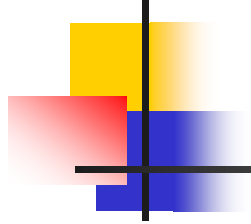


P. Allen



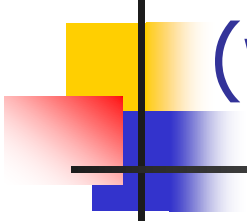
M. Rao

Well-quasi-order



Well-quasi-order

(well partial order = partial well-order)



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Definition. *A partial order on graphs is a well partial order if it contains no infinite antichains, i.e. no infinite sets of graphs pairwise incomparable with respect to the order.*

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Partial orders on graphs:

- induced subgraph relation
- subgraph relation
- minor relation
- induced minor relation
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[Robertson, Neil](#); [Seymour, P. D.](#) **Graph minors. XX. Wagner's conjecture.** *J. Combin. Theory Ser. B* 92 (2004), no. 2, 325–357

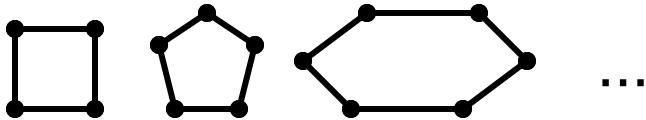


N. Robertson

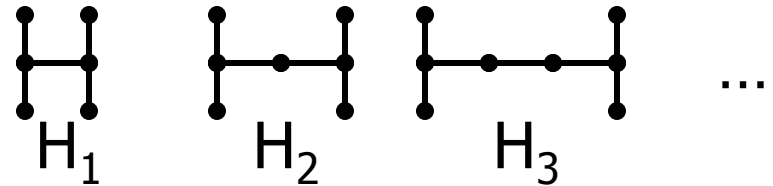
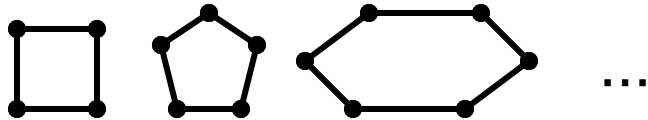


P. Seymour

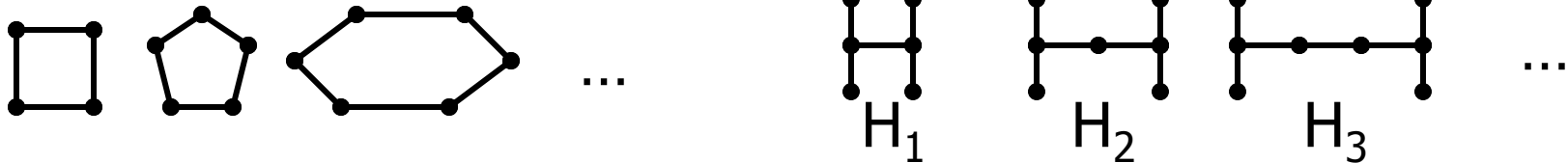
Infinite antichains with respect to subgraph and induced subgraph relations



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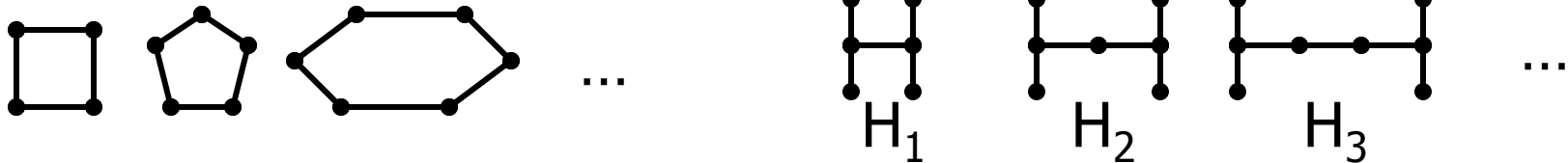
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Ding, Guoli **Subgraphs and well-quasi-ordering.**
J. Graph Theory 16 (1992) 489--502.

Theorem. *A class of graphs closed under taking subgraphs is wqo by the subgraph relation if and only if it contains finitely many cycles and finitely many H-graphs*

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Ding, Guoli

Position: Professor of Mathematics at LSU

Ph.D.: RUTCOR, Rutgers University 1991. Advisor P.D. Seymour



Hereditary classes of graphs well-quasi-ordered by induced subgraphs



Damaschke, Peter Induced subgraphs and well-quasi-ordering. *J. Graph Theory* 14 (1990), no. 4, 427–435.

Petkovšek, Marko Letter graphs and well-quasi-order by induced subgraphs. *Discrete Math.* 244 (2002), no. 1-3, 375–388.

Korpelainen, N., Lozin, V. Two forbidden induced subgraphs and well-quasi-ordering. *Discrete Math.* 311 (2011), 1813-1822

Korpelainen, N., Lozin, V. Bipartite induced subgraphs and well-quasi-ordering. *J. Graph Theory* 67 (2011), 235–249.



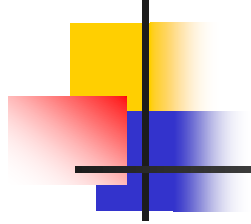
Well-quasi-orderability implies
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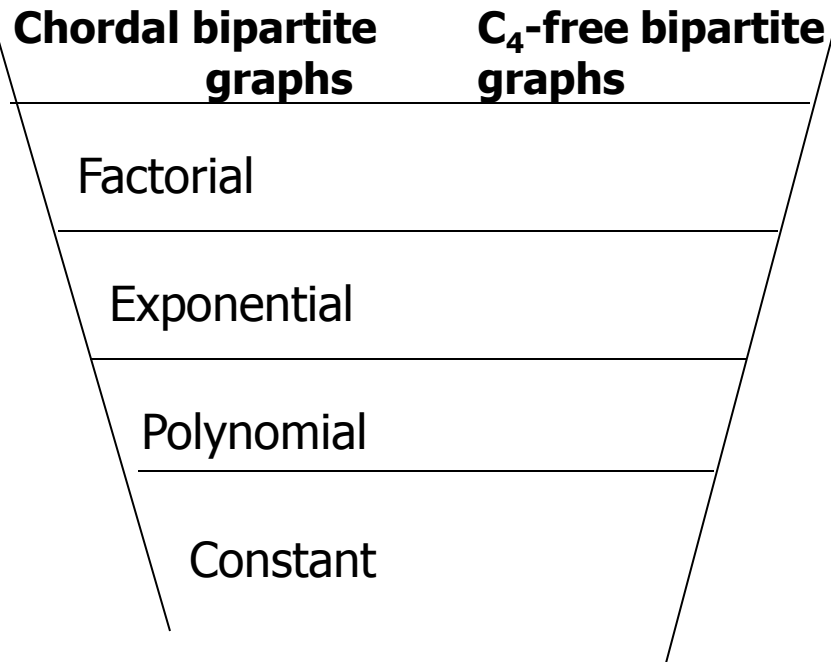
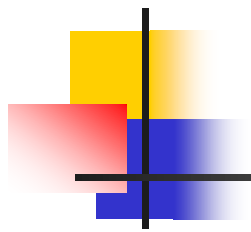


Daligault, Jean; Rao, Michael; Thomassé, Stéphan Well-quasi-order of relabel functions. *Order* 27 (2010), no. 3, 301–315,

Every well-quasi-ordered hereditary class is at most factorial?



Every well-quasi-ordered hereditary class is at most factorial?





Subclasses of chordal bipartite and C_4 -free bipartite graphs without a forest

Theorem. *For any forest F , the classes of F -free chordal bipartite graphs and (F, C_4) -free bipartite graphs are (at most) factorial.*



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Theorem. *For any forest F , the classes of F -free chordal bipartite graphs and (F, C_4) -free bipartite graphs are (at most) factorial.*

Corollary. *Every subclass of chordal bipartite or C_4 -free bipartite graphs which is wqo by the induced subgraph relation is (at most) factorial.*



Chordal bipartite graphs

(without a forest)

Let $G=(U,V,E)$ be a bipartite graph with bipartition U,V .

The bipartite adjacency matrix of G is a 0-1 matrix whose rows correspond to U and columns correspond to V .



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In a binary matrix, a Γ is a pair of rows and columns that induce the following submatrix:

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Fact. *A bipartite graph is chordal bipartite if and only if its bipartite adjacency matrix has a Γ -free ordering.*



Chordal bipartite graphs

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In a Γ -free matrix a 1 entry is called redundant if replacing it with 0 results in a Γ .



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Lemma. *Let F be a forest. Then the number of non-redundant 1s in a chordal bipartite F -free graph with n vertices is at most $O(n)$.*



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Proof (sketch). W.l.o.g. let F be a rooted tree in which every non-leaf vertex has exactly d children and every leaf vertex is of distance d from the root.



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Let $A(1)$ be the set of the first d non-redundant 1s taken from each row.

Let $B(1)$ be the set of the first d non-redundant 1s taken from each column which are not in $A(1)$.



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Let $A(i)$ be the set of the first d non-redundant 1s taken from each row which are not in $A(1) \cup B(1) \cup A(2) \cup B(2) \cup \dots \cup A(i-1) \cup B(i-1)$.

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Proof (sketch). $A(d+1)$ is empty.

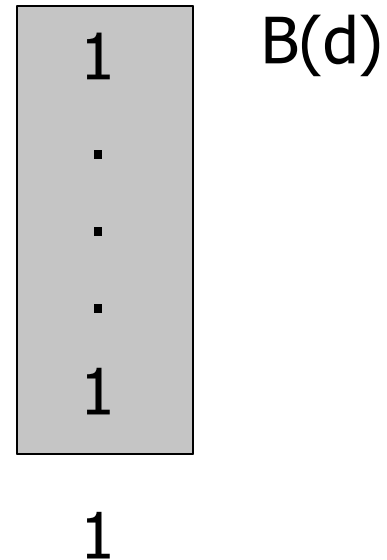
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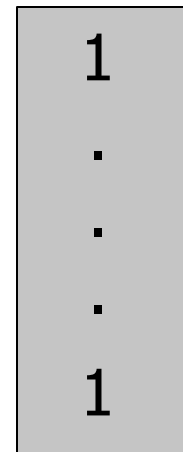
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A(d)



B(d)

1

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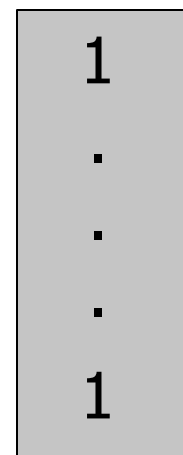
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$A(d)$



$B(d)$

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Therefore, the number of non-redundant 1s is at most $2dn$.



Chordal bipartite graphs

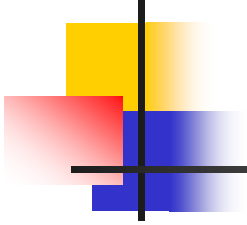
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Corollary. *Every F -free chordal bipartite graph can be represented with $O(n \log n)$ bits, and therefore, the class of F -free chordal bipartite graphs is at most factorial.*

Subclasses of C_4 -free bipartite graphs





Subclasses of C_4 -free bipartite graphs without long induced paths



Subclasses of C_4 -free bipartite graphs without long induced paths

Corollary. *For every k , the class of (P_k, C_4) -free bipartite graphs is well-quasi-ordered and is of bounded clique-width.*



Subclasses of C_4 -free bipartite graphs without long induced paths

Theorem. *For every k, n, p , the class of $(P_k K_n K_{p,p})$ -free graphs is well-quasi-ordered and is of bounded clique-width.*

Corollary. *For every k , the class of $(P_k C_4)$ -free bipartite graphs is well-quasi-ordered and is of bounded clique-width.*



Graphs containing long paths

Theorem. *Every graph containing a long path (not necessarily induced) contains either a big clique or a big induced biclique or a long induced path.*



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Corollary. *For every k, n, p , there exists an s such that the class of $(P_k, K_n, K_{p,p})$ -free graphs is a subclass of the class of graphs containing no P_s as a subgraph (not necessarily induced).*



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Fact. *Any class of graphs closed under taking subgraphs (not necessarily induced) and containing no long paths is of bounded tree-width and well-quasi-ordered.*



Be yourself no matter what they say

Sting

“Englishman in New York”



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Thank you